# On the chromatic number of set systems 

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#### Abstract

An $(r, l)$-system is an $r$-uniform hypergraph in which every set of $l$ vertices lies in at most one edge. Let $m_{k}(r, l)$ be the minimum number of edges in an $(r, l)$-system that is not $k$-colorable. Using probabilistic techniques, we prove that $$
a_{r, l}\left(k^{r-1} \ln k\right)^{\frac{l}{l-1}} \leq m_{k}(r, l) \leq b_{r, l}\left(k^{r-1} \ln k\right)^{\frac{l}{l-1}},
$$ where $b_{r, l}$ is explicitly defined and $a_{r, l}$ is sufficiently small. We also give a different argument proving (for even $k$ ) $$
m_{k}(r, l) \geq a_{r, l}^{\prime} k^{\frac{(r-1) l}{l-1}}
$$ where $a_{r, l}^{\prime}=\frac{r-l+1}{r}\left(2^{r-1} r e\right)^{\frac{-l}{l-1}}$. Our results complement earlier results of Erdős and Lovász [10] who mainly focused on the case $l=2, k$ fixed, and $r$ large.


## 1 Introduction

A hypergraph $H$ is $k$-colorable if its vertex set can be partitioned into $k$ color classes, such that no edge is monochromatic. The chromatic number $\chi(H)$ of $H$ is the minimum $k$ such

[^0]that $H$ is $k$-colorable. A classical extremal problem is to determine the minimum number of edges in an $r$-uniform hypergraph ( $r$-graph for short) that is not $k$-colorable. This minimum has been denoted $m_{k}(r)$ (see [2, 4, 6, 8, 9, 12, 14] for the results in the case $k=2$ and [1] for large $k$ ). If we restrict to the class of simple hypergraphs, i.e., those where every two distinct vertices lie in at most one edge, then the corresponding parameter is denoted by $m_{k}^{*}(r)$. This parameter was first studied by Erdős and Lovász [10]. They proved the bounds
\[

$$
\begin{equation*}
\frac{k^{2(r-2)}}{16 r(r-1)^{2}} \leq m_{k}^{*}(r) \leq 1600 r^{4} k^{2(r+1)} \tag{1}
\end{equation*}
$$

\]

which imply that

$$
\lim _{r \rightarrow \infty} m_{k}^{*}(r)^{1 / r}=k^{2}
$$

We consider a larger class of hypergraphs. A partial $(r, l)$-system (henceforth, $(r, l)$ system), is an $r$-uniform hypergraph in which every set of $l$ vertices lies in at most one edge. Let $m_{k}(r, l)$ be the minimum number of edges in an $(r, l)$-system that is not $k$-colorable; thus $m_{k}^{*}(r)=m_{k}(r, 2)$.

The works $[13,15,11]$ on Steiner systems with small independence number yield results for ( $r, l$ )-systems, and imply upper bounds on $m_{k}^{*}(r)$ which improve (1) for $k$ very large in comparison with $r$. In particular, Grable, Phelps and Rödl [11] for every $r$ and infinitely many $k$ constructed simple hypergraphs (in fact, Steiner systems) with chromatic number at least $k+1$ and at most $c 4^{r} r^{2} k^{2 r-2} \ln ^{2} k$ edges. Thus, for such $r$ and $k$,

$$
\begin{equation*}
m_{k}^{*}(r) \leq c 4^{r} r^{2} k^{2 r-2} \ln ^{2} k \tag{2}
\end{equation*}
$$

Our first result improves the upper bound in (1) in the range $k^{4}>0.01 r^{6}\left(\ln ^{2} e k\right)$. It has the advantage over (2) that it applies for every $l \geq 2$. We write $(r)_{l}$ for $r(r-1) \cdots(r-l+1)$.

Theorem 1 Let $r \geq 3, l \geq 2$. Then

$$
m_{k}(r, l) \leq 2 \frac{\left(c_{r, l}\right)^{l}}{(r)_{l}}\left(k^{r-1} \ln e k\right)^{\frac{l}{l-1}}
$$

where $c_{r, l}=\left(2 r^{3 l}\right)^{\frac{1}{l-1}}$.
We also improve the lower bound in (1) for $r \geq 3$ and large $k$.
Theorem 2 Let $r \geq 3$. If $k$ is even, then

$$
m_{k}(r, l) \geq d_{r, l} k^{\frac{(r-1) l}{l-1}}
$$

where

$$
d_{r, l}=\left[\frac{1}{\left(2^{r-1} r e\right)^{l}} \prod_{i=1}^{l-1}\left(1-\frac{i}{r}\right)\right]^{\frac{1}{l-1}}
$$

It is easy to see that this implies the result stated in the abstract. In the case where $r$ is fixed, we match the order of magnitude of the upper bound of Theorem 1.

Theorem 3 Let $r>l \geq 2$ be fixed. Then there exists $c$ depending only on $r$ and $l$ such that, for sufficiently large $k$ we have $m_{k}(r, l) \geq c\left(k^{r-1} \ln k\right)^{\frac{l}{l-1}}$.

We prove Theorem 1 in section 2 and Theorem 2 in section 3. In section 4 we generalize a result from [13] about the chromatic number of hypergraphs with large independent sets; this result is used in section 5 in the proof of Theorem 3.

## 2 The upper bound

The bounds of the kind (2) in $[13,15,11]$ hold for all $r$ and $k$, but apply only to large $k$ as written. Our construction also works for every $r>2$ and $k \geq 2$. It is an example of a random greedy algorithm.

Proof of Theorem 1: Consider the following procedure:
(1) order all $r$-element subsets of the set $\{1,2, \ldots, n\}$ at random: $R_{1}, \ldots, R_{\binom{n}{r}}$;
(2) Construct the family $G_{0}, \ldots, G_{\binom{n}{r}}$ of hypergraphs with the vertex set $V=\{1, \ldots, n\}$ as follows: $G_{0}$ has no edges and for $j=1, \ldots,\binom{n}{r}$ if $G_{j-1}+R_{j}$ is an $(r, l)$-system, then we let $G_{j}=G_{j-1}+R_{j}$, otherwise, $G_{j}=G_{j-1}$; (3) Let $G(n, r)=G_{\binom{n}{r}}$.

Clearly, Part (2) is a deterministic procedure once the ordering is defined. Our aim is to prove that if $n=\left\lceil c_{r, l}\left(k^{r-1} \ln e k\right)^{\frac{1}{l-1}}\right\rceil$, where $c_{r, l}=\left(2 r^{3 l}\right)^{\frac{1}{l-1}}$ then with positive probability $G(n, r)$ has no independent set of size $\lceil n / k\rceil$. Thus such a hypergraph has no $k$-colorings. Since $G(n, r)$ is an ( $r, l$ )-system by construction, this will give us (for $r \geq 3$ ) an example of an $(r, l)$-system with chromatic number at least $k+1$ and the number of edges at most

$$
\frac{\binom{n}{l}}{\binom{r}{l}} \leq \frac{n(n-1)^{l-1}}{(r)_{l}} \leq 2 \frac{c_{r, l}^{l}}{(r)_{l}}\left(k^{r-1} \ln e k\right)^{\frac{l}{l-1}}
$$

The proof follows from the following claim.
Claim. For an arbitrary set $X$ of vertices in $G(n, r)$ of cardinality $x=\lceil n / k\rceil$, the probability that $X$ induces no edges in $G(n, r)$ is less than $\binom{n}{x}^{-1}$.

Proof. Fix an $X$ of size $x=\lceil n / k\rceil$. Let $B_{X}$ be the event that $X$ induces no edges in $G(n, r)$. Observe that $B_{X}$ implies that every $r$-set $T \subseteq X$ must be preceded (in the random ordering) by some $r$-set $R$ not in $X$ such that $R \in G(n, r)$ and $|R \cap T| \geq l$. Consequently, $l \leq|R \cap X| \leq r-1$. Let us call such an $R$ a witness for $T \in[X]^{r}$ not being included in $G(n, r)$. The point is that if $B_{X}$ happens, then we must have a large number of witnesses in $G(n, r)$, and the probability of the latter is small. Indeed, each $R \in G(n, r)$ can be a witness for at most $\binom{r-1}{l}\binom{x-l}{r-l} r$-sets $T \subset X$. This means that in order to prevent all $\binom{x}{r} r$-sets $T$ of $X$ to appear in $G(n, r)$, the number of witnesses has to be at least

$$
m=\left\lceil\frac{\binom{x}{r}}{\binom{r-1}{l}\binom{x-l}{r-l}}\right\rceil=\left\lceil\frac{(x)_{l}}{(r)_{l}\binom{(-1}{l}}\right\rceil .
$$

For $j \geq 1$, let $A_{j}=A_{X, j}$ denote the event that the first $j$ edges $R_{l_{1}}, R_{l_{2}}, \ldots, R_{l_{j}}$ in $G(n, r)$ such that $\left|R_{l_{i}} \cap X\right| \geq l$ are not contained in $X$, i.e. $l \leq\left|R_{l_{i}} \cap X\right| \leq r-1$. The previous paragraph yields that if $B_{X}$ occurs, then $A_{m}$ also occurs.

The rest of the proof consists of bounding the probability of $A_{m}$ from above by $\binom{n}{x}^{-1}$.
For this calculation, we further assume that $R_{l_{1}}$ is the witness that appears first in the ordering, and that for each $1<j \leq m, R_{l_{j}}$ is the first witness which comes after $R_{l_{j-1}}$. Let $G^{j}=G_{l_{j}-1}$ be the family of all $r$-sets included in $G(n, r)$ before the $j$ th witness $R_{l_{j}}$ is chosen. For $1 \leq j \leq m$, let $\mathcal{S}_{j}$ be the collection of all $r$-sets $S$, such that $|X \cap S| \geq l$ and $|R \cap S|<l$ for all $R \in G^{j}$.

Since $A_{m} \subset A_{m-1} \subset \cdots \subset A_{2} \subset A_{1}$, we have

$$
\mathbf{P}\left\{A_{m}\right\}=\mathbf{P}\left\{A_{1}\right\} \cdot \mathbf{P}\left\{A_{2} \mid A_{1}\right\} \cdot \ldots \cdot \mathbf{P}\left\{A_{m} \mid A_{m-1}\right\}
$$

To estimate these probabilities we first note that each of the events $A_{1}$ and $A_{j+1} \mid A_{j}, j=$ $1, \ldots, m-1$ corresponds to a random choice from the set $\mathcal{S}_{j}$ with the result that the chosen set belongs to $\mathcal{S}_{j}-[X]^{r}$.

Since $\left|\mathcal{S}_{1}\right| \leq\binom{ x}{l}\binom{n}{r-l}$ we have

$$
\mathbf{P}\left\{A_{1}\right\}=\frac{\left|\mathcal{S}_{1}\right|-\binom{x}{r}}{\left|\mathcal{S}_{1}\right|} \leq 1-\frac{\binom{x}{r}}{\binom{x}{l}\binom{n}{r-l}} .
$$

Furthermore, suppose that $j>1$, and let $j \leq m_{0}=\left\lceil\frac{m}{2}\right\rceil$. Assume now that the event $A_{j}$ occurred. Since

$$
\mathbf{P}\left\{A_{j} \mid A_{j-1}\right\}=\frac{\left|\mathcal{S}_{j}-[X]^{r}\right|}{\left|\mathcal{S}_{j}\right|}=1-\frac{\left|\mathcal{S}_{j} \cap[X]^{r}\right|}{\left|\mathcal{S}_{j}\right|}
$$

we need to estimate the cardinality of the set $\mathcal{S}_{j} \cap[X]^{r}$.

The hypergraph $G^{j}$ contains precisely $j-1 r$-sets $R$ with $l \leq|R \cap X| \leq r-1$. Each of these is a witness for at most

$$
\binom{|X \cap R|}{l}\binom{x-l}{r-l} \leq\binom{ r-1}{l}\binom{x-l}{r-l}
$$

$r$-sets. Consequently, the number of $r$-sets $T$ in $X$ with no witness at this stage is

$$
\left|\mathcal{S}_{j} \cap[X]^{r}\right| \geq\binom{ x}{r}-(j-1)\binom{r-1}{l}\binom{x-l}{r-l} \geq\binom{ x}{r}-\left(m_{0}-1\right)\binom{r-1}{l}\binom{x-l}{r-l} \geq \frac{1}{2}\binom{x}{r}
$$

where the last inequality follows from the choice of $m_{0}$. Summarizing, we infer that

$$
\mathbf{P}\left\{A_{j} \mid A_{j-1}\right\} \leq 1-\frac{\frac{1}{2}\binom{x}{r}}{\binom{n}{r-l}\binom{x}{l}}
$$

This yields

$$
\begin{gathered}
\mathbf{P}\left\{A_{m}\right\} \leq \mathbf{P}\left\{A_{m_{0}}\right\}=\mathbf{P}\left\{A_{1}\right\} \cdot \mathbf{P}\left\{A_{2} \mid A_{1}\right\} \cdot \ldots \cdot \mathbf{P}\left\{A_{m_{0}} \mid A_{m_{0}-1}\right\} \leq \\
\leq\left(1-\frac{\frac{1}{2}\binom{x}{r}}{\binom{n}{r-l}\binom{x}{l}}\right)^{m_{0}} \leq\left(1-\frac{\frac{1}{2}\binom{x}{r}}{\binom{n}{r-l}\binom{x}{l}}\right)^{\frac{(x)_{l}}{2(r)_{l}\binom{r-1}{l}}} \leq \exp \left\{-\frac{(x)_{r}}{4 n^{r-l}\binom{r}{l}(r)_{l}\binom{(-1}{l}}\right\} .
\end{gathered}
$$

In order to prove the claim we will show that the last expression is less than $\binom{n}{x}^{-1}$. Since $\binom{n}{x}<(n e / x)^{x}=\exp \left\{x \ln \frac{e n}{x}\right\}$, and $x=\lceil n / k\rceil$, we have

$$
\mathbf{P}\left\{A_{m}\right\} \cdot\binom{n}{x}<\exp \left\{x\left(\ln \frac{e n}{x}-\frac{(x-1)_{r-1}}{4 n^{r-l}\binom{r}{l}(r)_{l}\binom{r-1}{l}}\right)\right\} .
$$

By the choice of $n$, it is easy to observe that $x \geq n / k \geq(r-1)^{3}$. Thus for $r \geq 3$

$$
(x-1)_{r-1} \geq\left(1-\frac{r-1}{x}\right)^{r-1} x^{r-1} \geq\left(1-\frac{1}{(r-1)^{2}}\right)^{r-1} x^{r-1} \geq \frac{1}{2} x^{r-1} \geq \frac{1}{2}\left(\frac{n}{k}\right)^{r-1}
$$

Consequently, $\mathbf{P}\left\{A_{m}\right\} \cdot\binom{n}{x}$ is strictly less than

$$
\exp \left\{x\left(\ln e k-\frac{\frac{1}{2}\binom{n}{k}^{r-1}}{4 n^{r-l}\binom{r}{l}(r)_{l}\binom{r-1}{l}}\right)\right\} \leq 1
$$

where the last inequality follows since

$$
n \geq\left(2 r^{3 l}\right)^{\frac{1}{l-1}}\left(k^{r-1} \ln e k\right)^{\frac{1}{l-1}} \geq\left[8\binom{r}{l}(r)_{l}\binom{r-1}{l}\right]^{\frac{1}{l-1}}\left(k^{r-1} \ln e k\right)^{\frac{1}{l-1}}
$$

## 3 Lower bounds from the Lovász Local Lemma

In this section, we prove Theorem 2. Our main tool is the symmetric version of the Lovász Local Lemma which we state below (see [5] for a proof).

Lemma 4 (Local Lemma) Let $A_{1}, \ldots, A_{n}$ be events in a probability space. Suppose that each event $A_{i}$ is mutually independent of a set of all the other events $A_{j}$ but at most $d$, and that $\operatorname{Prob}\left[A_{i}\right] \leq p$ for all $i$. If $e p(d+1) \leq 1$, then $\operatorname{Prob}\left[\wedge \overline{A_{i}}\right]>0$.

We use the following lemma from [10], whose proof we supply for completeness.

Lemma 5 Let $H$ be an r-graph. If every vertex in $H$ has degree at most $k^{r-1} / e r$, then $\chi(H) \leq k$.

Proof. We color the vertices of $H$ with $k$ colors, with each color being assigned to each vertex independently with equal probability. The probability that a given edge is monochromatic is $1 / k^{r-1}$. The event $A_{F}$ that edge $F$ is monochromatic is independent of all events $A_{F^{\prime}}$ with $F \cap F^{\prime}=\emptyset$. The number of $F^{\prime}$ with $F \cap F^{\prime} \neq \emptyset$ is at most $r\left(k^{r-1} /(e r)-1\right) \leq k^{r-1} / e-1$. The Local Lemma (Lemma 4) therefore implies that there is a $k$-coloring with no monochromatic edge.

Proof of Theorem 2: Let $H$ be an $(r, l)$-system with at most $z=c_{r, l} k^{\frac{(r-1) l}{l-1}}$ edges, where

$$
c_{r, l}=\left[\frac{1}{\left(2^{r-1} r e\right)^{l}} \prod_{i=1}^{l-1}\left(1-\frac{i}{r}\right)\right]^{\frac{1}{l-1}} .
$$

Let

$$
A=\left\{v \in V(H): \operatorname{deg}(v)>\frac{k^{r-1}}{e r 2^{r-1}}\right\}
$$

Let $B=V(H)-A$, and let $H_{A}$ and $H_{B}$ be the subhypergraphs induced by $A$ and $B$, respectively. By Lemma 5 , there is a proper $k / 2$-coloring of $H_{B}$.

Now color $H_{A}$ randomly using a new set of $k / 2$ colors, where each color appears on each vertex independently with equal probability. Since $H$ is an $(r, l)$-system, every vertex in $H_{A}$ has degree (in $H_{A}$ ) at most

$$
\Delta=\frac{(a-1)(a-2) \cdots(a-l+1)}{(r-1)(r-2) \cdots(r-l+1)}
$$

where $a=|A|$. Consequently, each edge $E$ in $H_{A}$ is incident with at most $d=\Delta r-r$ other edges. Moreover, since $H$ has at most $z$ edges and $z r>a k^{r-1} /\left(e r 2^{r-1}\right)$, we infer that $a<z e r^{2} 2^{r-1} / k^{r-1}$. Consider the space of colorings with each vertex being colored randomly and independently of others. For each edge $E$ in $H_{A}$, let $M_{E}$ be the event that $E$ is monochromatic. Since $p=\mathbf{P}\left\{M_{E}\right\}=(2 / k)^{r-1}$ and

$$
\begin{aligned}
e p(d+1) & \leq e\left(\frac{2}{k}\right)^{r-1}\left(\frac{(a-1)(a-2) \cdots(a-l+1)}{(r-1)(r-2) \cdots(r-l+1)} r-r+1\right) \leq \\
& \leq e\left(\frac{2}{k}\right)^{r-1} \frac{a^{l-1} r}{(r-1)(r-2) \cdots(r-l+1)} \leq \\
& \leq e^{l} \frac{r^{2 l-1}}{(r-1) \cdots(r-l+1)}\left(\frac{2}{k}\right)^{(r-1) l} z^{l-1}<1,
\end{aligned}
$$

the Local Lemma implies that there is a proper $k / 2$-coloring of $H_{A}$. These two colorings together yield a proper $k$-coloring of $H$.

## 4 From independent sets to proper colorings

In this section, we prove a preliminary Lemma 8 to our main lower bound, Theorem 3, which might be interesting of its own. A special case appears in [13]. The following fact was kindly pointed out to us by a referee.

Lemma 6 Let $f(m)$ be a monotonically non-decreasing function, $f(1)=1$, and $f(m) \leq m$ for every $m$. Let $G$ be a graph on $n$ vertices. Let $I_{1}, \ldots, I_{t}$ be a family of disjoint independent sets in $G$ with $i_{l}=\left|I_{l}\right|$ for $l=1, \ldots, t$. Let $x_{0}=0$ and $x_{l}=\sum_{j=1}^{l} i_{j}$. If $i_{j} \geq f\left(n-x_{j-1}\right)$ for every $j=1, \ldots, t$, then $t \leq \sum_{l=n-x_{t}+1}^{n} \frac{1}{f(l)}$.

Proof. Since $f(m)$ is monotonically non-decreasing and $i_{j} \geq f\left(n-x_{j-1}\right)$, we have

$$
\sum_{l=n-x_{t}+1}^{n} \frac{1}{f(l)}=\sum_{j=1}^{t} \sum_{l=n-x_{j}+1}^{n-x_{j-1}} \frac{1}{f(l)} \geq \sum_{j=1}^{t} \sum_{l=n-x_{j}+1}^{n-x_{j-1}} \frac{1}{f\left(n-x_{j-1}\right)}=\sum_{j=1}^{t} \frac{i_{j}}{f\left(n-x_{j-1}\right)} \geq t
$$

This lemma (due to a referee) directly implies the following nice corollary.
Lemma 7 Let $f(m)$ be a monotonically non-decreasing function, $f(1)=1$, and $f(m) \leq m$ for every $m$. Let $G$ be a graph on $n$ vertices. If for every $2 \leq m \leq n$, the independence number of every $m$-vertex subgraph of $G$ is at least $f(m)$, then $\chi(G) \leq \sum_{j=1}^{n} \frac{1}{f(j)}$.

Lemma 8 Let $0 \leq \alpha<1$ and $\beta<1-\alpha$. Let $H$ be a hypergraph with $n$ vertices. Suppose that every subhypergraph $P$ of $H$ (including $H$ itself) with $m \geq 2$ vertices has an independent set of size at least $f(m)=c m^{\alpha}(\ln e m)^{\beta}$ for some constant $c>0$. Then there is a $d=$ $d(c, \alpha, \beta)>0$ such that $\chi(H) \leq d n^{1-\alpha}(\ln e n)^{-\beta}$.

Proof. Define $f(1)=1$. Then by Lemma 7,

$$
\begin{gathered}
\chi(H) \leq \sum_{j=1}^{n} \frac{1}{f(j)} \leq 1+\int_{1}^{n} \frac{1}{c} x^{-\alpha}(\ln e x)^{-\beta} d x \leq \\
\leq 1+\frac{1}{c(1-\alpha-\beta)} \int_{1}^{n} x^{-\alpha}(\ln e x)^{-\beta}\left(1-\alpha-\frac{\beta}{\ln e x}\right) d x= \\
=1+\left.\frac{1}{c(1-\alpha-\beta)}\right|_{1} ^{n} x^{1-\alpha}(\ln e x)^{-\beta}=1+\frac{1}{c(1-\alpha-\beta)}\left(n^{1-\alpha}(\ln e n)^{-\beta}-1\right) .
\end{gathered}
$$

This proves the lemma.
We use Lemma 8 to prove that $(r, l)$-systems with not too many vertices are $k$-colorable. The following result of Rödl and Šinajová guarantees large independent sets in such $r$-graphs.

Theorem 9 (Rödl-S̆inajová [15]) Let $H$ be an ( $r, l$ )-system on $n$ vertices. Then $H$ has an independent set of size at least $c n^{(r-l) /(r-1)}(\ln n)^{1 /(r-1)}$, where $c$ is a positive constant depending only on $r$ and $l$.

Theorem 9 together with Lemma 8 implies
Theorem 10 Let $H$ be an $(r, l)$-system on $n$ vertices. Then $\chi(H) \leq c\left(n^{l-1} / \ln n\right)^{1 /(r-1)}$ for some constant $c$ depending only on $r$ and $l$. Moreover, there is another constant $c^{\prime}$ (also depending only on $r$ and $l$ ) such that, if $n \leq c^{\prime}\left(k^{r-1} \ln k\right)^{1 /(l-1)}$, then $\chi(H) \leq k$.

## 5 The main lower bound

In this section we prove Theorem 3. The main idea to properly $k$-color the $(r, l)$-system is to greedily take maximal independent sets. We therefore need a lower bound on the size of a maximal independent set in an $r$-graph. Such a bound is provided for a fairly restricted class by a result of Ajtai et al. [3].

Theorem 11 ([3]) Let $G$ be an r-uniform hypergraph without 2-,3- and 4-cycles. If $|E(G)| /|V(G)|$ is very large in comparison with $r$, then

$$
\begin{equation*}
\alpha(G) \geq c \frac{|V(G)|^{r /(r-1)}}{|E(G)|^{1 /(r-1)}}\left(\ln \frac{|E(G)|}{|V(G)|}\right)^{1 /(r-1)} \tag{3}
\end{equation*}
$$

where $c$ depends only on $r$.

We remark that the condition $|E(G)| /|V(G)|$ being large can be removed by changing the constant $c$. Duke et al. [7] extended this bound (with a different constant) to the class of simple hypergraphs. We need the following generalization of [7] for $(r, l)$-systems; the proof follows from the idea in [15].

Theorem 12 Let $r, l$ be integers with $r>l>1$, and let $\delta=(r-l) /(8 r-10)$. Suppose that $F$ is an $(r, l)$-system with $|V(F)|=n$ and $|E(F)| \geq n^{l-\delta}$. Then

$$
\begin{equation*}
\alpha(F) \geq c_{1} n\left(\frac{\ln w}{w}\right)^{1 /(r-1)} \tag{4}
\end{equation*}
$$

where $w=|E(F)| / n$ and $c_{1}$ depends only on $r$ and $l$.

Proof. (Sketch) Let $\epsilon_{0}=(4 l-5) /(4 r-5)$ and $\epsilon_{1}=(l-1-\delta) /(r-1)<(l-1) /(r-1)$. Set $\epsilon=\left(\epsilon_{0}+\epsilon_{1}\right) / 2$. Consider a random induced subsystem $H$ of $F$, where every vertex in $H$ is included with probability $p=n^{-\epsilon}$ independently of all other vertices. The expected number of vertices in $H$ is $p n$ and the expected number of edges in $H$ is $p^{r}|E(F)|$. In [15] it is proven that with positive probability, we can delete at most half of the vertices of $H$ to obtain a subsystem $G$ of $F$ with
(1) no cycles of length less than five,
(2) $\mid V(G)=p n / 2$, and
(3) $|E(G)| \leq 2 p^{r}|E(F)|$.

We apply Theorem 11 to $G$. With given $|V(G)|$, the bound we seek for $\alpha(G)$ decreases when $|E(G)|$ grows. Therefore, letting $z=|E(F)|$, we obtain that $\alpha(F)$ is at least

$$
\alpha(G) \geq c \frac{(p n / 2)^{\frac{r}{r-1}}}{\left(2 p^{r} z\right)^{\frac{1}{r-1}}}\left(\ln \frac{4 p^{r} z}{p n}\right)^{\frac{1}{r-1}}=c \frac{(n / 2)^{\frac{r}{r-1}}}{(2 z)^{\frac{1}{r-1}}}\left(\ln \frac{4 z}{n^{\epsilon(r-1)+1}}\right)^{\frac{1}{r-1}} \geq c_{1} \frac{n^{\frac{r}{r-1}}}{z^{\frac{1}{r-1}}}\left(\ln \frac{z}{n}\right)^{\frac{1}{r-1}}
$$

The last inequality follows by replacing the exponent inside the logarithm by a factor outside of the logarithm.

Proof of Theorem 3: Let $H$ be an $(r, l)$-system with at most $c_{2}\left(k^{r-1} \ln k\right)^{\frac{l}{l-1}}$ edges. Set $c_{3}=1 /\left(e r 3^{r-1}\right)$. Partition $V(H)$ into two parts:
$V_{0}$ - vertices of degree at most $c_{3} k^{r-1}$, and
$V_{1}$ - vertices of degree greater than $c_{3} k^{r-1}$.
By Lemma 5 there is a proper coloring of the vertices in $V_{0}$ with $k / 3$ colors. It remains to properly color the vertices of $V_{1}$ with at most $2 k / 3$ colors.

Let $H_{1}=H\left(V_{1}\right)$ and $n_{1}=\left|V_{1}\right|$. Let $c_{4}$ be chosen so that by Theorem 10, every $(r, l)$ system with at most

$$
\begin{equation*}
n_{0}=c_{4}\left(k^{r-1} \ln k\right)^{1 /(l-1)} \tag{5}
\end{equation*}
$$

vertices is $\frac{k}{3}$-colorable. Because

$$
n_{1} c_{3} k^{r-1} \leq \sum_{v \in V_{1}} \operatorname{deg}(v) \leq r|E(H)|
$$

we obtain

$$
\begin{equation*}
n_{1} \leq \frac{r c_{2}}{c_{3}}\left(k^{r-1} \ln ^{l} k\right)^{1 /(l-1)}=\frac{r c_{2}}{c_{3} c_{4}} n_{0} \ln k . \tag{6}
\end{equation*}
$$

Let $a_{1}=\frac{n_{1}}{n_{0}}$. If $a_{1} \leq 1$, then we are done, and by (6), in any case,

$$
\begin{equation*}
a_{1} \leq \frac{r c_{2}}{c_{3} c_{4}} \ln k \tag{7}
\end{equation*}
$$

Let $i \geq 1$ and consider the following procedure:

## Step $i$ :

(a) If $a_{i}>1$, then we distinguish between two cases depending on whether $\left|E\left(H_{i}\right)\right|$ is large.
(I) If $\left|E\left(H_{i}\right)\right| \geq n_{i}^{l-\delta}$, where $\delta$ is as in Theorem 12 , then we can apply Theorem 12. Choose in $H_{i}$ a maximum independent set $I_{i}$, let $H_{i+1}=H_{i}-I_{i}, n_{i+1}=\left|V\left(H_{i+1}\right)\right|=n_{i}-\left|I_{i}\right|$ and $a_{i+1}=\frac{n_{i+1}}{n_{0}}$. Now go to Step $i+1$.
(II) If $\left|E\left(H_{i}\right)\right|<n_{i}^{l-\delta}$, then partition $V\left(H_{i}\right)$ into two sets $X$ and $Y$, where $X$ consists of all vertices of $H_{i}$ with degree less than $d k^{r-1}$, with $d=1 /\left(\right.$ er $\left.^{r-1}\right)$.

By the choice of $d$, Lemma 5 implies that the hypergraph induced by $X$ can be properly $k / 6$-colored. Let $d^{\prime}$ be chosen so that, by Theorem 10 , every $(r, l)$-system on at most $d^{\prime}\left(k^{r-1} \ln k\right)^{1 /(l-1)}$ vertices is properly $k / 6$-colorable. Because

$$
|Y| d k^{r-1} \leq \sum_{v \in Y} \operatorname{deg}(v) \leq r\left|E\left(H_{i}\right)\right| \leq r n_{i}^{l-\delta}<r n_{1}^{l-\delta}
$$

we conclude that since $k$ is sufficiently large

$$
\begin{equation*}
|Y| \leq \frac{r n_{1}^{l-\delta}}{d k^{r-1}} \leq \frac{r}{d k^{r-1}}\left[\frac{r c_{2}}{c_{3}} \ln k\left(k^{r-1} \ln k\right)^{\frac{1}{l-1}}\right]^{l-\delta} \leq d^{\prime}\left(k^{r-1} \ln k\right)^{\frac{1}{l-1}} \tag{8}
\end{equation*}
$$

Consequently, the subhypergraph of $H_{i}$ induced by $Y$ can be properly $k / 6$-colored. These two colorings together yield a proper $k / 3$-coloring of $H_{i}$. Color $H_{i}$ properly with $k / 3$ colors. Since all vertices of $H$ are now colored, we stop the procedure.
(b) if $a_{i} \leq 1$, then the number of vertices in the uncolored hypergraph is at most $n_{0}$. We apply Theorem 10 to color these vertices with $k / 3$ colors. Now we stop the procedure.

Suppose that the procedure stops on Step $t+1$. That means that in Steps $i=1, \ldots, t$, we were in Case (a) part (I). We will prove that $t \leq k / 3$. Observe that this implies that $H$ is $k$-colorable.
$(\alpha)$ We used $k / 3$ colors to color $H\left(V_{0}\right)$,
$(\beta)$ We used $t$ colors for $I_{1}, \ldots, I_{t}$,
$(\gamma)$ Regardless of whether we stopped the procedure due to Case (a) part (II), or Case (b), in each situation we used $k / 3$ new colors. This yields the required $k$-coloring of $H$.

In order to complete the argument, we will show that in $(\beta)$ we have $t \leq k / 3$.
By the definition of $a_{i}$, we have

$$
\left|E\left(H_{i}\right)\right| /\left|V\left(H_{i}\right)\right| \leq \frac{c_{2}\left(k^{r-1} \ln k\right)^{l /(l-1)}}{a_{i} c_{4}\left(k^{r-1} \ln k\right)^{1 /(l-1)}} \leq \frac{c_{2}\left(k^{r-1} \ln k\right)^{l /(l-1)}}{a_{i} c_{4}\left(k^{r-1} \ln k\right)^{1 /(l-1)}}=\frac{c_{2} k^{r-1} \ln k}{a_{i} c_{4}}
$$

and hence by Theorem 12 for large $k$,

$$
\begin{equation*}
\left|I_{i}\right| \geq c_{1} n_{i}\left(\frac{a_{i} c_{4}}{c_{2} k^{r-1} \ln k}\right)^{\frac{1}{r-1}}\left(\ln \frac{c_{2} k^{r-1} \ln k}{a_{i} c_{4}}\right)^{\frac{1}{r-1}} \geq\left(\frac{a_{i} c_{4}}{c_{2}}\right)^{\frac{1}{r-1}} \frac{c_{1} n_{i}}{k}=\left(\frac{c_{4}}{c_{2} n_{0}}\right)^{\frac{1}{r-1}} \frac{c_{1} n_{i}^{\frac{r}{r-1}}}{k} \tag{9}
\end{equation*}
$$

Let $c_{5}=c_{1}\left(\frac{c_{4}}{c_{2}}\right)^{\frac{1}{r-1}}$. Then by (9), the conditions of Lemma 6 are satisfied with $f(m)=$ $\frac{c_{5} \frac{r}{\frac{r}{r-1}}}{k n_{0}^{\frac{1}{r-1}}}$. Hence by Lemma 6,

$$
\begin{gathered}
t \leq 1+\sum_{l=n_{0}+1}^{n} \frac{1}{f(l)}=1+\frac{k n_{0}^{\frac{1}{r-1}}}{c_{5}} \sum_{l=n_{0}+1}^{n} l^{-\frac{r}{r-1}} \leq 1+\frac{k n_{0}^{\frac{1}{r-1}}}{c_{5}} \int_{n_{0}}^{n} x^{-\frac{r}{r-1}} d x \leq \\
\leq 1+\frac{k(r-1) n_{0}^{\frac{1}{r-1}}}{c_{5}}\left(n_{0}^{\frac{-1}{r-1}}-n^{\frac{-1}{r-1}}\right) \leq 1+\frac{k(r-1)}{c_{5}} .
\end{gathered}
$$

Thus if we choose $c_{2}$ small enough to make $c_{5}>6(r-1)$, then for large $k$ we will have $t<k / 3$. This proves the bound.

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