

CHAPTER 9: RAMSEY THEORY

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ABSTRACT. Ramsey's theorem asserts that if the k -element subsets of a large set are colored with a small number of colors, then there is a homogeneous (monochromatic) set H all of whose k -element subsets are assigned the same color. This important generalization of the Pigeon Hole Principle implies that complete disorder is impossible. Instead, buried inside large systems are subsystems with complete uniformity.

1. BASIC NOTATION AND TERMINOLOGY

For a positive integer n , we let $[n] = \{1, 2, \dots, n\}$. For a set X and a non-negative integer k , $\binom{X}{k}$ denotes the set of all k -element subsets of X . In Ramsey theoretic settings, it is common to refer to a map $\phi : \binom{X}{k} \rightarrow R$ as a *coloring* of the k -element subsets of X , and the elements of R will be referred to as *colors*. Typically, we will just take $R = [r]$ for some positive integer r , so ϕ will also be called an r -coloring.

When $\phi : \binom{X}{k} \rightarrow [r]$ is a r -coloring, a subset $H \subseteq X$ is called a *homogeneous set* (also a *monochromatic set*) when there exists a color $\alpha \in [r]$ so that $\phi(A) = \alpha$ for every $A \in \binom{H}{k}$.

Theorem 1 (Ramsey's theorem). *If k and r are positive integers, and (h_1, h_2, \dots, h_r) are integers with $h_i \geq k$ for $i = 1, 2, \dots, r$, then there exists a least positive integer $t_0 = R(k : h_1, h_2, \dots, h_r)$ so that if X is any set with $|X| \geq t_0$, then for every r -coloring $\phi : \binom{X}{k} \rightarrow [r]$ of the k -element subsets of X , there exists an $\alpha \in [r]$ and a subset $H \subseteq X$ with $|H| \geq h_i$ so that $\phi(A) = \alpha$ for every $A \in \binom{H}{k}$.*

Proof. We use a double induction. The first induction is on k , and the second is on r . When $k = 1$, the result holds for all r as the theorem reduces to a restatement of the Pigeonhole Principle. As an aside, we note that

$$R(1 : h_1, h_2, \dots, h_r) = 1 + \sum_{i=1}^r h_i - 1.$$

However, in general, we will not be able to say much about the exact value of Ramsey numbers. Instead, the emphasis is on the fact that they *exist*!

Now assume validity for some $k \geq 1$ and consider the next value of k . Now the induction is on r . When $r = 1$, it is easy to see that $R(k : h_1) = h_1$. Now consider the case $r = 2$.

Let $q = R(1; h_1, h_2)$ and define a sequence of numbers $(s_0, s_1, s_2, \dots, s_q)$ as follows. First set $s_0 = k - 1$. If s_i has been defined, and $1 \leq i < q$, set $s_{i+1} = 1 + R(k - 1, s_i, s_i)$. Note that $s_1 = 1 + R(k - 1; k - 1, k - 1) = 1 + (k - 1) = k$. We show that $R(k : h_1, h_2)$ exists and it is at most s_q .

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For additional data, do a web search and look for Stanley Radziszowski, who maintains the most current information on his web site.

3. ESTIMATING RAMSEY NUMBERS

We will find it convenient to utilize the following approximation due to Stirling. You can find a proof in almost any advanced calculus book.

$$n! \equiv \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \left(1 + \frac{1}{12n} + \frac{1}{288n^2} - \frac{139}{51840n^3} + O\left(\frac{1}{n^4}\right)\right).$$

Of course, we will normally be satisfied with the first term:

$$n! \equiv \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$$

Using Stirling's approximation, we have the following upper bound:

$$R(2; n, n) \leq \binom{2n-2}{n-1} \equiv \frac{2^{2n}}{4\sqrt{\pi n}}$$

Here is an exponential lower bound.

Theorem 3.

$$R(2; n, n) \geq (1 + o(1)) \frac{n}{e\sqrt{2}} 2^{\frac{1}{2}n}$$

Proof. Let t be an integer with $t > n$ and consider the following probability space. The outcomes in the probability space are graphs with vertex set $\{1, 2, \dots, t\}$. For each i and j with $1 \leq i < j \leq t$, edge ij is present in the graph with probability $1/2$. Furthermore, the events for distinct pairs are independent.

Let X_1 denote the random variable which counts the number of n -element subsets of $\{1, 2, \dots, t\}$ for which all $\binom{n}{2}$ pairs are edges in the graph. Similarly, X_2 is the random variable which counts the number of n -element subsets of $\{1, 2, \dots, t\}$ for which all $\binom{n}{2}$ pairs are edges are *not* in the graph. Then set $X = X_1 + X_2$.

By linearity of expectation, $E(X) = E(X_1) + E(X_2)$ while

$$E(X_1) = E(X_2) = \binom{t}{n} \frac{1}{2^{\binom{n}{2}}}.$$

If $E(X) < 1$, then there must exist a graph with vertex set $\{1, 2, \dots, t\}$ without a K_n or an I_n . We then consider the inequality

$$2 \binom{t}{n} \frac{1}{2^{\binom{n}{2}}} < 1$$

Using the Stirling approximation, we see that this inequality holds when

$$t \geq (1 + o(1)) \frac{n}{e\sqrt{2}} 2^{\frac{1}{2}n}$$

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